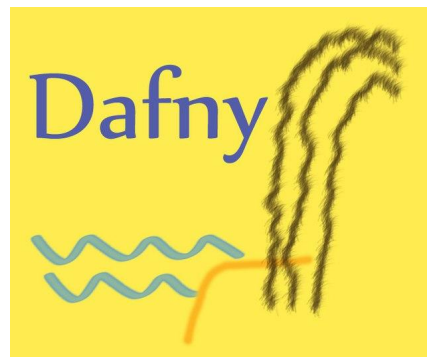


Metamorph: Synthesizing Large Objects with Dafny APIs

Aleksandr Fedchin, **Alexander Bai**, Jeff Foster

Dafny

- Verification-aware language that allows writing specification in first order logic.
- Check the program w.r.t. the specification with SMT solvers
- UNSAT means verified, a satisfying assignment will give back an counterexample



Dafny is Used for Critical Systems

AWS Encryption SDK &
Authorization Engine



Dafny EVM



Project Everest



Ironclad/Ironfleet



Motivation

- Existing tools mostly restrict themselves to functional code
- Many synthesis tools have a max number of methods
- Synthesis tools usually require execute candidate solutions
- Dafny allows **full** specification of API methods, including of side-effects

Motivating Example: Social Network API

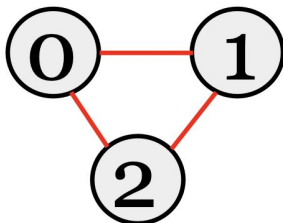
```
datatype User = User(name:nat,  
    friends:set<nat>,  
    incomingRequests:set<nat>)  
  
class SocialNetwork {  
    ghost var users:map<nat,User>  
  
    constructor()...  
    method AddUser(id:nat)...  
    method RequestConnect(from:nat, to:nat) ...  
    method RemoveUser(id:nat)...  
    method RequestResponse(from:nat, to:nat, response:bool)...  
}
```

Method Specification Example

```
method AddUser(id:nat)
  requires id !in users
  modifies this
  ensures users == old(users)[id := User(id, {}, {})]
{
  users := users[id := User(id, {}, {})]
}
```

Example Synthesis Query

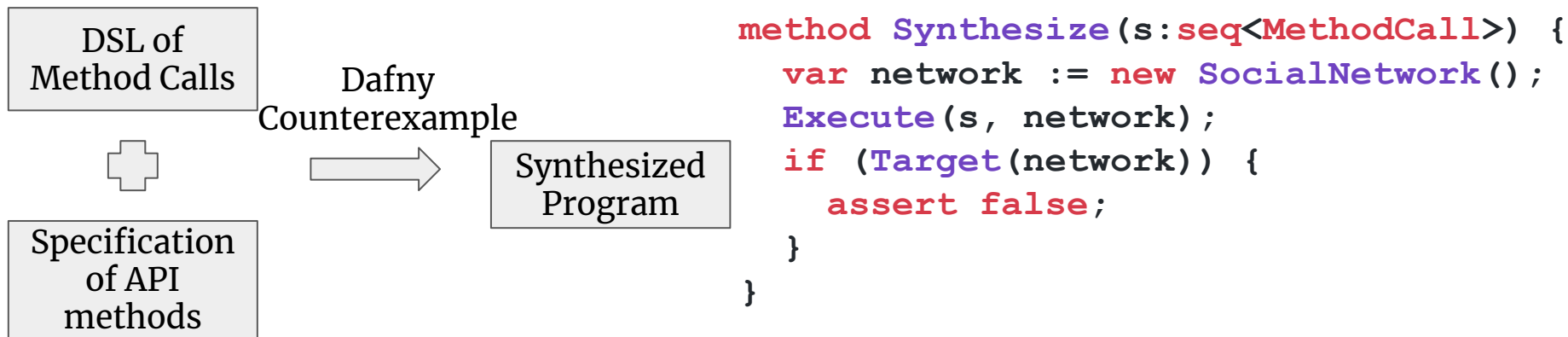
```
predicate Target(s:SocialNetwork)
  reads db
{
  s.users.Keys == {0, 1, 2} &&
  0 in s.users[1].friends &&
  1 in s.users[2].friends &&
  2 in s.users[0].friends
}
```



```
static method Solution()
  returns (s:SocialNetwork)
  ensures Target(result)
{
  s := new SocialNetwork();
  s.AddUser(1);
  s.AddUser(0);
  s.RequestConnect(1, 0);
  s.RequestResponse(0, 1, true);
  s.AddUser(2);
  s.RequestConnect(1, 2);
  s.RequestResponse(2, 1, true);
  s.RequestConnect(2, 0);
  s.RequestResponse(0, 2, true);
}
```

Solution 1: Using Dafny Directly

We can use Dafny counterexamples to synthesize the goal directly.



Problem: The excessive use of quantifiers makes this approach impractical

What if we reason about one method at a time?

Backward Synthesis by Counterexample

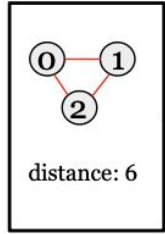
- Consider the following assertion:

```
method InferPreviousState(s: SocialNetwork, arg: nat)
  modifies receiver {
    assume !Target(s);
    assume arg in s.users;
    s.RemoveUser(arg);
    assert !Target(s);
  }
```

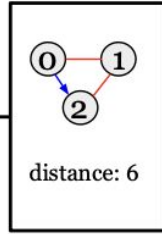
- The counterexample will give us the the previous state as a conjunction of constraints

Backward Synthesis as Graph Search

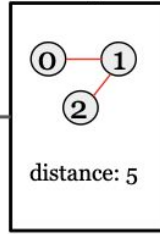
Target State



S₁

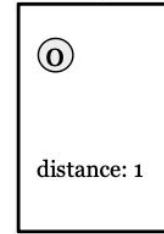


S₃

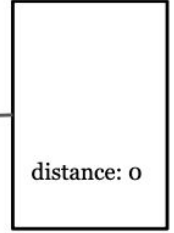


...

S₇



Initial State



RequesResponse
(2, 0, true)

RequestConnect
(0, 2)

RequesResponse
(2, 1, true)

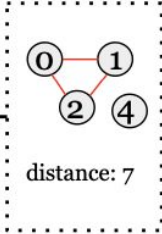
RemoveUser(4)

RemoveUser(4)

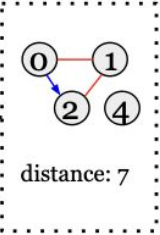
AddUser(0)

RemoveUser(4)

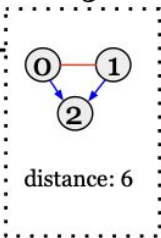
S₂



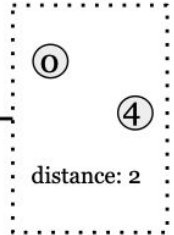
S₆



S₅



S₈



Distance Metric

- Counterexamples are variable assignment, so they represent object states as conjunctions of constraints.
- We abstract away concrete values in these constraints...

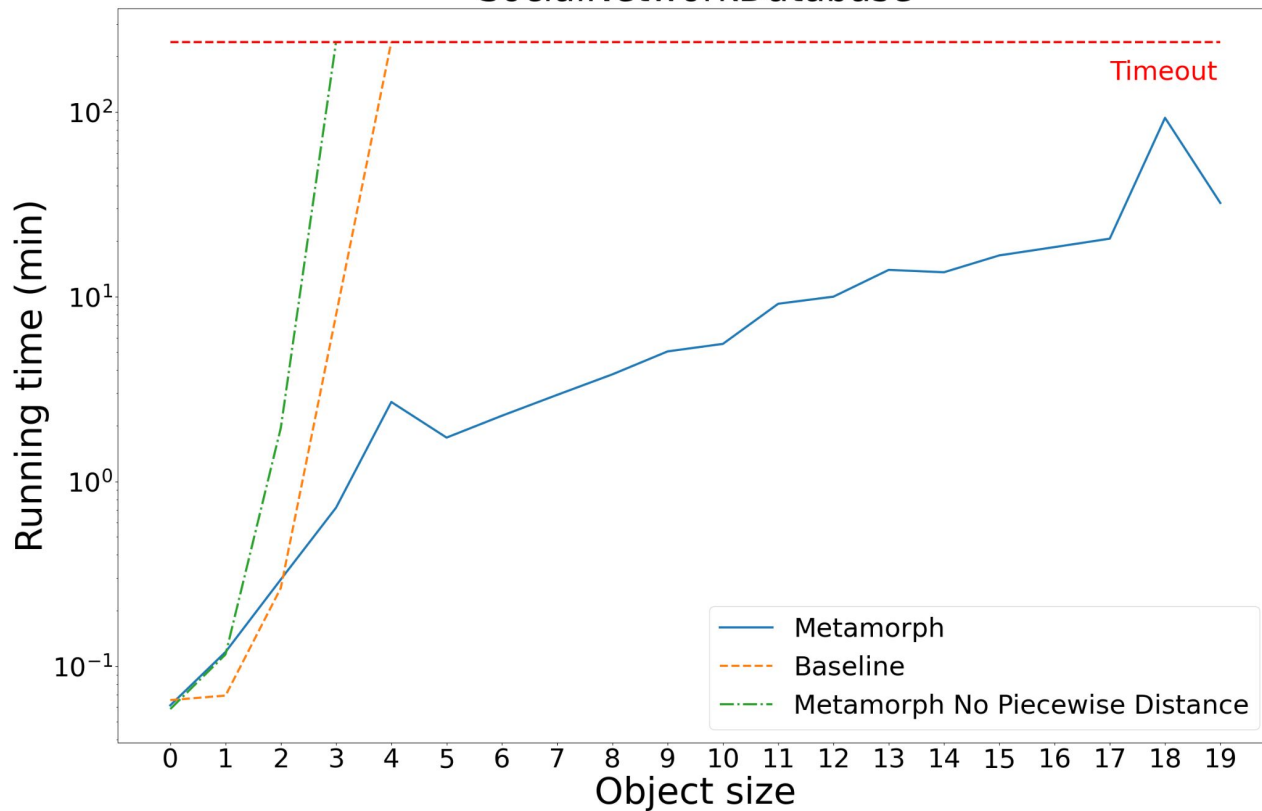
`0 in users && 1 in users <=> X(a) && X(b) && a != b`

- ...and reason about how each API method interacts with each symbolic constraint.

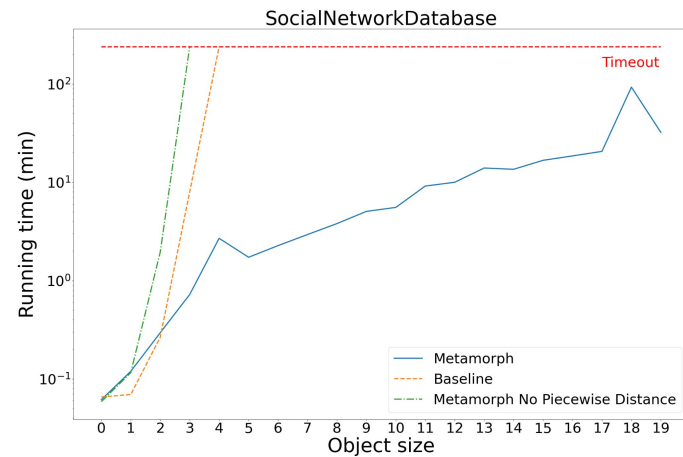
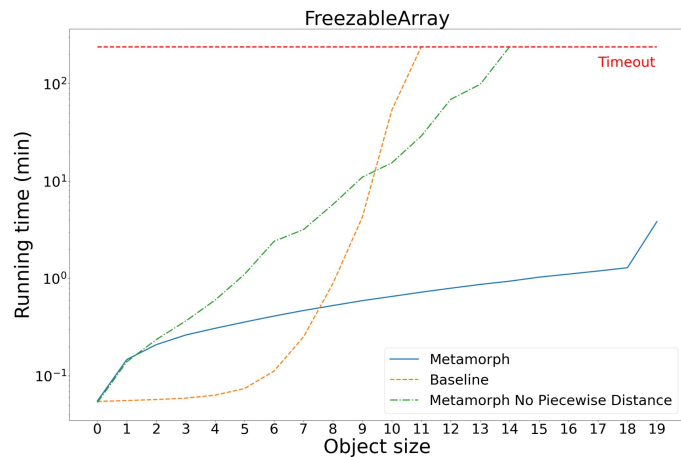
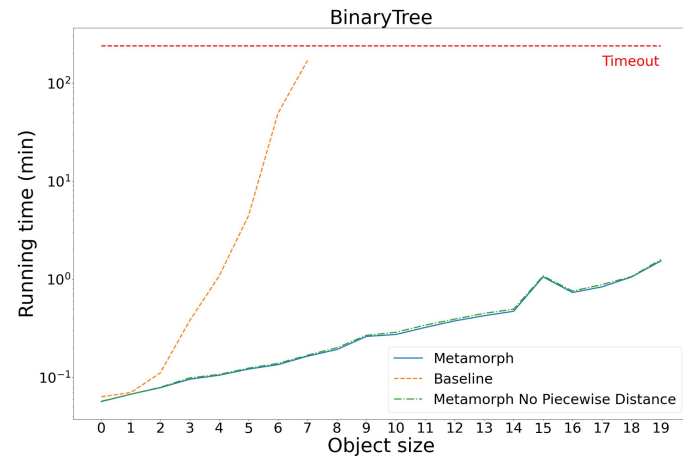
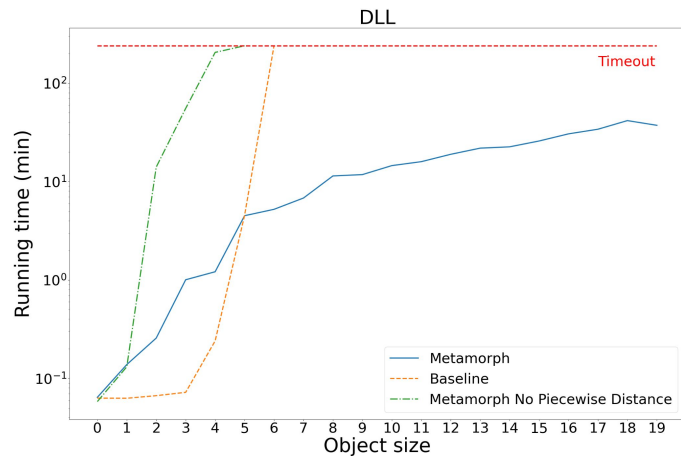
`X(a) -> RemoveUser(a) -> ¬X(a)`

Evaluation

SocialNetworkDatabase



Evaluation



Future Work

- Integrate with Dafny Test Generation framework.
- Synthesizing programs with non-linear control flow.
- Combining forward and backward search

Summary

- We introduced an approach of using counterexamples to “concretize” states
- We used a distance metric to guide the backward synthesis

